ECE 273 Convex Optimization and Applications Instructor: Jun-Kun Wang Scribe: Merlin Chang, Marialena Sfyraki May 21, 2024 Editors/TAs: Marialena Sfyraki

Lecture 15: Min-Max Optimization

1 Introduction

This lecture and next lecture concern solving

$$
\inf_{x \in X} \sup_{y \in Y} g(x, y).
$$

In general, we want to use inf and sup rather than min and max because the minimum or maximum might not exist. Every time we know for sure the inf*,*sup are attained, we can substitute them with min*,* max.

1.1 Applications

There are a few interpretations of the min-max optimization problems.

1.1.1 Two Player Game

One immediate intuitive understanding of this problem would be in the form of a two-player game, in which *y* wants to maximize $g(\cdot, \cdot)$, while *x* would like to minimize $g(\cdot, \cdot)$.

1.1.2 Empirical Risk Minimization

Recall that Empirical risk minimization is the following,

$$
\min_{x} \sum_{i=1}^{n} f_i(x)
$$

One way we can more robustly formulate ERM is to consider it as a 2 player zero-sum game. Our goal is try to minimize our worst loss. Which is

$$
\min_{x} \max_{i \in [n]} f_i(x) \iff \min_{x} \max_{y \in \Delta_n} y_i f_i(x)
$$

This is an example of "distributionally robust optimization" problems.

1.1.3 Bi-linear Game

Let's consider the following setting:

$$
\min_{x \in \Delta_n} \max_{y \in \Delta_m} y^\top A x
$$

where $\Delta_n := \{ x \in \mathbb{R}^n : \sum_{i=1}^n x[i] = 1, x[i] \ge 0 \}$. We can see that

$$
y^{\top} Ax = \sum_{i,j} y[i] x[j] A_{i,j}
$$

=
$$
\sum_{i,j} Pr(i_y = i) Pr(j_x = j) A_{i_y, j_x}
$$

=
$$
\mathbb{E}_{i_y \sim P_y, j_x \sim P_x} [A_{i_y, j_x}]
$$

Since *x* is trying to minimize and *y* is trying to maximize, we can view the above optimization *y [⊤]Ax* is the loss for *x* and reward for *y* when they choose their action from their respective distributions.

Furthermore, if we let *A* be the mis-classification matrix as

$$
A_{i,j} = \begin{cases} 1 & , \text{ if } h_i(z_j) \neq l_j \\ 0 & , otherwise \end{cases}
$$

where we let *z* be a point in the training set $\{z_j \in \mathbb{R}^d, l_j = \{+1, -1\}\}_{j=1}^m$ and $H :=$ ${h_i(\cdot)}_{i=1}^n$ be a set of prediction functions, i.e., $h_i(\cdot) : \mathbb{R}^d \to \{+1, -1\}$, the min max problem becomes

$$
\min_{x \in \Delta_n} \max_{y \in \Delta_m} y^\top A x \equiv \min_{x \in \Delta_n} \max_{y \in \Delta_m} \sum_{i,j} y[i] x[j] 1[h_i(z_j) \neq l_j].
$$

In this case, *x* is trying to minimize loss, which can interpreted as searching for the best prediction function. *y* is trying to maximize reward, which is same as finding the hardest point to to predict adversarially.

2 Theory

We need to introduce a few notions about what does it mean to solve a min-max problem.

Definition 1 (Saddle points/ Nash Equilibrium). Let $x \in X$ and $y \in Y$ and $g(\cdot, \cdot)$: $X \times Y \to \mathbb{R}$ *. A pair of points* $(x_*, y_*) \in X \times Y$ *is a saddle point of* $g(\cdot, \cdot)$ *if*

$$
g(x_*, y) \le g(x_*, y_*) \le g(x, y_*)
$$

for any $x \in X$ *and* $y \in Y$.

Remark 1: When *x* and *y* are at equilibrium (x_*, y_*) , they have no incentive to change. Given x_* , if *y* switches from y_* to any point $y \in \mathcal{Y}$, the value of $g(\cdot, \cdot)$ will not be greater than that of $g(x_*, y_*)$. Similarly, given y_* , if *x* switches from x_* to any point $x \in \mathcal{X}$, the value of $g(\cdot, \cdot)$ will not be smaller than that of $g(x_*, y_*)$.

Remark 2: The goal of solving a min-max optimization problem is to find the saddle point.

The following lemma is always true.

Lemma 1. Let $g(\cdot, \cdot): X \times Y \to \mathbb{R}$, where *X* and *Y* are not empty. Then,

$$
\inf_{x \in X} \sup_{y \in Y} g(x, y) \ge \sup_{y \in Y} \inf_{x \in X} g(x, y).
$$

Proof. By the definition of infimum, we know that for any $x' \in X$

$$
g(x', y) \ge \inf_{x \in X} g(x, y)
$$

This also implies that

$$
\sup_{y \in Y} g(x', y) \ge \sup_{y \in Y} \inf_{x \in X} g(x, y)
$$

for any $x' \in X$, thus we know that the x that achieves the infimum also satisfy the above inequality, thus

$$
\inf_{x \in X} \sup_{y \in Y} g(x, y) \ge \sup_{y \in Y} \inf_{x \in X} g(x, y).
$$

Theorem 1. Let $g(x, y) : X \times Y \to \mathbb{R}$, where X and Y are not empty. A point (*x∗, y∗*) *is a saddle point if and only if*

- The supremum in $\sup_{y\in Y} \inf_{x\in X} g(x, y)$ is attained at y_* & the infimum in $\inf_{x\in X} \sup_{y\in Y} g(x, y)$ *is attained at x∗.*
- *Also,* sup_{*y*∈*Y*} inf_{*x*∈*X*} $g(x, y) = \inf_{x \in X} \sup_{y \in Y} g(x, y)$ *, lemma 1 becomes equality.*

Proof. We will show both directions.

• (\Longrightarrow) : Let (x_*, y_*) be a saddle point for *g*. We know by definition of the saddle point

$$
g(x_*, y) \le g(x_*, y_*) \le g(x, y_*)
$$

for any $x \in X$ and $y \in Y$, we have

$$
\inf_{x \in X} \sup_{y \in Y} g(x, y) \le \sup_{y \in Y} g(x_*, y) = g(x_*, y_*) = \inf_{x \in X} g(x, y_*) \le \sup_{y \in Y} \inf_{x \in X} g(x, y) \tag{1}
$$

Thus

$$
\inf_{x \in X} \sup_{y \in Y} g(x, y) \le \sup_{y \in Y} \inf_{x \in X} g(x, y).
$$

By lemma 1, we know that the above inequality is strict equality, where

$$
\inf_{x \in X} \sup_{y \in Y} g(x, y) = \sup_{y \in Y} \inf_{x \in X} g(x, y).
$$

This equality also implies that the following in [\(1](#page-3-0)) are true,

$$
\inf_{x \in X} \sup_{y \in Y} g(x, y) = \sup_{y \in Y} g(x_*, y), \inf_{x \in X} g(x, y_*) = \sup_{y \in Y} \inf_{x \in X} g(x, y)
$$

which implies that infimum and supremum for $\inf_{x \in X} \sup_{y \in Y} g(x, y)$ and $\sup_{y \in Y} \inf_{x \in X} g(x, y)$ are attained at x_* and y_* respectively. This shows that if (x_*, y_*) is a saddle point, then the two items are satisfied.

 $\bullet~~(\Longleftarrow):$ Assume that we have the above two items in the theorem satisfied. We know by definition of sup and inf that

$$
\inf_{x \in X} g(x, y_*) \le g(x_*, y_*) \le \sup_{y \in Y} g(x_*, y) \tag{2}
$$

By the assumed conditions, we have

$$
\inf_{x \in X} g(x, y_*) = \sup_{y \in Y} \inf_{x \in X} g(x, y) = \inf_{x \in X} \sup_{y \in Y} g(x, y) = \sup_{y \in Y} g(x_*, y). \tag{3}
$$

Hence, by (2) (2) and (3) (3) , we have

$$
\inf_{x \in X} g(x, y_*) = g(x_*, y_*) = \sup_{y \in Y} g(x_*, y)
$$

which implies that if $g(x_*, y_*)$ is the supremum of g in terms of y, then

$$
g(x_*, y_*) \ge g(x_*, y), \forall y \in Y
$$

The same argument applies for *x*,

$$
g(x_*, y_*) \le g(x, y_*), \forall x \in X
$$

Thus, we achieved the condition for (x_*, y_*) to be a saddle point, where

$$
g(x_*,y) \le g(x_*,y_*) \le g(x,y_*), \forall x \in X, \forall y \in Y
$$

There are a few important implications of the above theorem.

Remark. If the inf sup problem and sup inf problem have different values, then there is no saddle point.

Remark. If a saddle point exists, then there might be multiple ones, and all of them must have the same minimax value, i.e., $\sup_{y \in Y} \inf_{x \in X} g(x, y) = \inf_{x \in X} \sup_{y \in Y} g(x, y)$ **Remark.** The theorem also implies that, if saddle points exists, the set of saddle points is the Cartesian product of $X^* \times Y^*$ when nonempty. The set X^*, Y^* are the set of solutions for the infimum and supremum of the following problems,

$$
\inf_{\mathbf{x}\in\mathbf{X}}\sup_{y\in Y}g(x,y),\ \sup_{\mathbf{y}\in\mathbf{Y}}\inf_{x\in X}g(x,y)
$$

This implies that the way we choose x_*, y_* is completely independent from each other.[\[1](#page-6-0)]

Let's look at a few examples to illustrate the theorem. **Example 1.** Let $g(x, y) = (x - y)^2, X = [-1, 1], Y = [-1, 1].$ From the *x*'s perspective,

$$
\inf_{x \in X} \sup_{y \in Y} (x - y)^2 = \inf_{x \in X} (1 + |x|)^2 = 1
$$

and from the *y*'s perspective,

$$
\sup_{y \in Y} \inf_{x \in X} (x - y)^2 = \sup_{y \in Y} 0 = 0
$$

thus, there is no saddle point.

Example 2. Let $g(x, y) = xy, X = (0, 1], Y = (0, 1].$ In this case, although the min-max optimization from both player's perspective is the same,

$$
\inf_{x \in X} \sup_{y \in Y} xy = \inf_{x \in X} x = 0 = \sup_{y \in Y} \inf_{x \in X} xy = \sup_{y \in Y} 0
$$

but in condition inf_{$x \in X$} $x = 0$, the infimum is not attained within *X*, thus saddle point still does not exist.

Theorem 1 tells us the general direction of finding a saddle point is the approach from both directions. Namely, we need to find a minimizer of *x* in $\sup_{y \in Y} g(x, y)$ and a maximizer of *y* in $\inf_{x \in X} g(x, y)$.

There are a few immediate ideas to measure progress to wards the saddle point, but they are have issues.

- Difference: $g(x, y) g(x_*, y_*)$. This value can be negative or equal to zero for an infinite number of points that are not saddle points.
- Norm distance: $||x x_*||_2^2 + ||y y_*||_2^2$. This quantity can go to zero at an arbitrarily slow rate

Let's consider splitting the problem into two according to theorem 1, where from the *x*'s perspective, let's try to solve

$$
\inf_{x \in X} \ell(x)
$$

where $\ell(x) := \sup_{y \in Y} g(x, y)$. In this view, we can use the standard measure of suboptimality gap, which is

$$
\ell(x) - \inf_{x \in X} \ell(x) = \sup_{y \in Y} g(x, y) - \inf_{x \in X} \sup_{y \in Y} g(x, y)
$$

The same idea can be applied to the *y*'s perspective and arrive at a similar measure of convergence when we only focus on the variable *y*,

$$
\sup_{y \in Y} h(y) - h(y) = \sup_{y \in Y} \inf_{x \in X} g(x, y) - \inf_{x \in X} g(x, y)
$$

where $h(y) := \inf_{x \in X} g(x, y)$. Finally, we can simply combine the above two measures to form a joint measure for saddle point progress,

$$
\sup_{y \in Y} g(x, y) - \inf_{x \in X} \sup_{y \in Y} g(x, y) + \sup_{y \in Y} \inf_{x \in X} g(x, y) - \inf_{x \in X} g(x, y)
$$

if we assume the existence of a saddle point, then the above measure reduces to

$$
\sup_{y \in Y} g(x, y) - \inf_{x \in X} g(x, y)
$$

since $\inf_{x \in X} \sup_{y \in Y} g(x, y) = \sup_{y \in Y} \inf_{x \in X} g(x, y)$ when saddle points exists.

Definition 2 (Duality Gap). For a function $g: X \times Y \to \mathbb{R}$, the duality gap on a *pair of points* $(x', y') \in X \times Y$ *is defined as*

$$
\sup_{y \in Y} g(x', y) - \inf_{x \in X} g(x, y')
$$

Remark. The duality gap is always non-negative even if saddle points don't exist. As $\sup_{y\in Y} g(x', y) \ge \inf_{x\in X} g(x, y')\forall x' \in X, \forall y' \in Y$. We will prove this property in the next lecture.

Bibliographic notes

The materials of this lecture is based on Chapter 11 of [[1\]](#page-6-0). Please refer to Chapter 11 of [[1](#page-6-0)] for more details about saddle points and its related theory.

References

[1] Francesco Orabona A Modern Introduction to Online Learning arXiv:1912.13213. 2023